

Ontology-based data integration in EPNet

Alessandro Mosca

KRDB Research Centre, Free University of Bozen-Bolzano, Italy

Joint work with D. Calvanese (UNIBZ), J. Remesal (UB), M. Rezk (GOOGLE), and G. Rull (SIRIS)

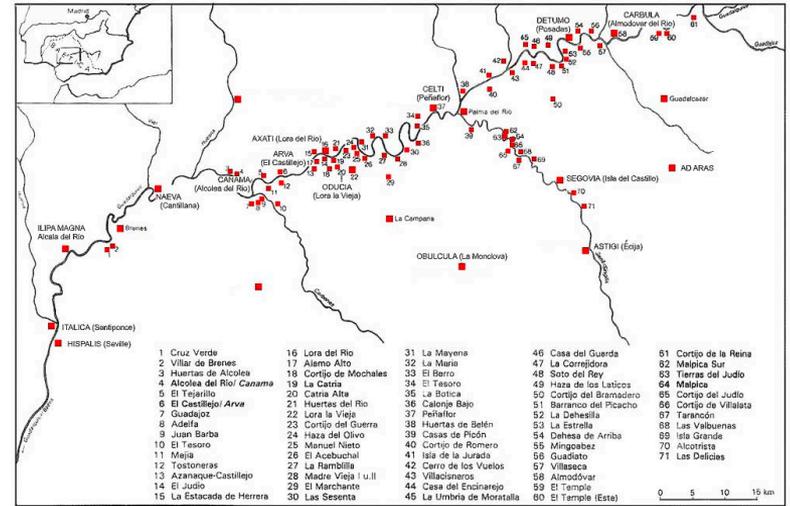
The Roman Empire trade system

The first complex European trade network

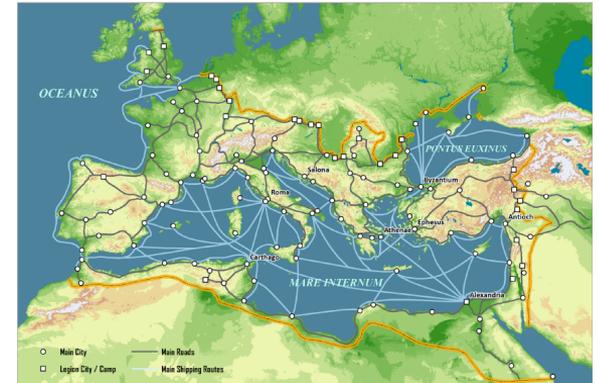


Tabula Peutingeriana
©Austrian National Library

“An integrated network of interactions and interdependencies between the Mediterranean basin and northern Europe”



Dressel 20 Kiln Sites @Lower Guadalquivir Valley (Remesal, 1997)



The Roman Empire transport system (c. 125 AD)

<http://people.hofstra.edu/>

The economy of the Roman Empire

An ongoing debate!

- ❑ Traditionally, the study of the food distribution during the Roman Empire was focused on the life of the city of Rome and on the long distance trade. More recently, the food distribution to the Roman army also gained a primary role in the overall understanding of the Roman Empire trade system.
- ❑ It has been argued that the emperor and his circle managed the relationship between food and army in order to supervise and control the whole Roman territory and to strengthen and maintain their own political power.
- ❑ Periphery and regional food distribution then obtain the status of first-class citizens in the understanding of the whole system.
- ❑ Starting from this innovative idea, different models have been proposed.



[House of Julia Felix](#) @Pompeii

Informal approaches are too informal...



Production and Distribution of Food during the Roman Empire:
Economic and Political Dynamics

- A plethora of unfalsifiable theories has been developed over the last couple of centuries to explain the organisation of the Roman Empire trade system.
- The ongoing debate remains exclusively speculative and often based on rhetoric.
- Specialists from History and Archaeology do not even consider the possibility that their hypotheses can be formally expressed as **theories that can be falsified⁽¹⁾ by analytical and computational methods.**

⁽¹⁾ Popper, Karl, *The Logic of Scientific Discovery*,
Basic Books, New York, NY, 1959

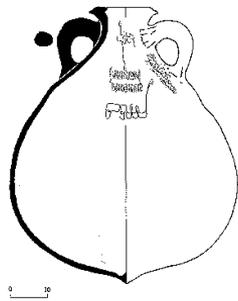
- A formal study of the mechanisms that have characterised the economic and political relations informing the Roman Empire trade system was still missing.



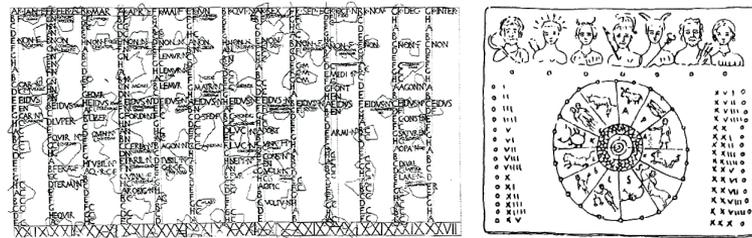
The data landscape



Edible items: olive oil, wine, etc.



Containers



Temporal context: production, distribution, etc.



Production places



Inscriptions: stamps and *tituli picti*



Distribution target places



Producers and Traders

Multiple datasets, physically distributed



number 1	
HD no.	↗ HD000165 (newer version at: ↗ EDR)
work status	checked with photo
province / Italic region	Etruria (Regio VII)
modern country	Italy
ancient find spot	Falerii (Novi), bei
modern find spot	
find spot (village, street, etc.)	
year	231 AD – 300 AD
literature ⓘ	AE 1982, 0274. (B) AE 1983, 0069. (B) I. Di Stefano Manzella, SupplIt 1 (Roma 1981) 143-144, Nr. 21; Foto. – AE 1982. L. Polverini, in: G. Barbieri (Hrsg.), Il lapidario Zeri di Mentana (Roma 1982) 104-105, Nr. 46; tav. 42. – AE 1983.
Transcription	D(is) M(anibus) s(acrum) / Aurelio Saturnino / militi torquato legionis / primes(II) Italic(a)je qui vix(it)!) an(n)is XL messibus(!) V1 / diebus X millitavit / annis XIII / Aurelius Arborius / fratri bene mer(enti) fecit

Baetis/Certis (river)

a Pleiades place resource

Creators: Jr., F.H. Stanley, R.C. Knapp
Contributors: R. Warner, R. Talbert, Sean Gillies, Tom Elliott, Jeffrey Becker
Copyright © The Contributors. Sharing and remixing permitted under terms of the Creative Commons Attribution 3.0 License (cc-by).
Last modified Jun 13, 2019 11:35 AM – History

The **Guadalquivir** river is the fifth longest of the Iberian peninsula.

Canonical URI for this page:
<https://pleiades.stoa.org/places/256010>

Representative Point (Latitude, Longitude):
37.868207, -2.910335

- Locations:**
- Barrington Atlas location (750 BC - AD 640)
 - location of course of Baetis/Certis fl. (modern)

- Names:**
- Baetis (750 BC - AD 640)
 - Certis (AD 300 - AD 640)
 - Perkes (unspecified date range)

- Baetis/Certis (river) makes connections with:**
- Baetis/Certis (river) ⇒ connection ⇒ Atlanticus Oceanus (unspecified date range)

- Baetis/Certis (river) receives connections from:**
- Unnamed Bridge ⇒ on ⇒ Baetis/Certis (river) (30 BC – AD 300)
 - Ligustinus L. ⇒ connection ⇒ Baetis/Certis (river) (unspecified date r.)
 - Guadaira (river) ⇒ connection ⇒ Baetis/Certis (river) (modern)
 - Singilis (river) ⇒ connection ⇒ Baetis/Certis (river) (unspecified date r.)
 - Ilipa Magna ⇒ on ⇒ Baetis/Certis (river) (30 BC - AD 300)
 - Anatolikon Stoma ⇒ connection ⇒ Baetis/Certis (river) (unspecified c.)
 - Baetis Aestuaria ⇒ connection ⇒ Baetis/Certis (river) (unspecified d.)

Place type:
river, drainage



Dressel 20

[previous amphora type](#) [next amphora type](#)

[details](#) | [characteristics](#) | [pictures](#) | [drawings](#) | [petrology](#) | [specimens](#) | [bibliography](#) | [3D models](#)

Variants of Dressel 20:

[\[Dressel 20 similis - Oliva 3\]](#) [\[Gauloise 14\]](#)

Distinctive Features

This has a large globular body with thick, sharply bent or oval shaped handles; short neck often with an internally concave rim and a small basal knob. This form clearly developed from the Augustan prototype Oberarden 83 and eventually evolved into the smaller Tejarillo 1. A broad scheme for the evolution of the Dressel 20 rim has been suggested by Berni (1998), provided under 'Drawings'. This form is commonly stamped towards the summit of the handle, with the name in relief set in a rectangular frame.

This is usually an abbreviated name of an individual, presumably the estate owner although place-names are also encountered (Remesal Rodríguez, 1986). It also includes complex painted inscriptions indicating amphora weights and individuals responsible for transport and control (Rodríguez Almeida, 1984). From the third century there is evidence of imperial ownership of some of these estates.

[See characteristics](#)

Date Range

The typical globular form was introduced by at least the Tiberian period (Xanten) and became established by the Claudian period (Colis et alii, 1977). Production continued up to the second half of the third century (Zevi, 1967; Blázquez Martínez & Remesal Rodríguez, 1999; 2001; 2003) (Monte Testaccio). In the western provinces it is the most common amphora from the late first to the early third centuries AD (Williams & Peacock, 1983; Remesal Rodríguez, 1986).

Search: [\[1st century AD\]](#) [\[2nd century AD\]](#) [\[3rd century AD\]](#)

Origin

Along the banks of the river Guadalquivir and its tributaries between Seville and Córdoba in the Roman province of *Baetica* (Clark-Maxwell, 1899; Bonser, 1931; Ponsich, 1974, 1979, 1991; Chic, 2001) in southern Spain, where kiln sites have been discovered ([Click](#) to map in Introduction). A graffito from Carlisle reads *ESVRI* (Vila Real de São Antonio) on the Algarve of Portugal (Mann, 1955), but the fabric is identical to the Spanish material. The vessel was imitated in Hispania *Tarraconensis*, along the coastal strip of *Baetica* and in *Germania* (Baudoux, 1996; Ehmig, 2003).

Search: [\[North West Europe\]](#) [\[Spain\]](#) [\[Western Mediterranean\]](#)



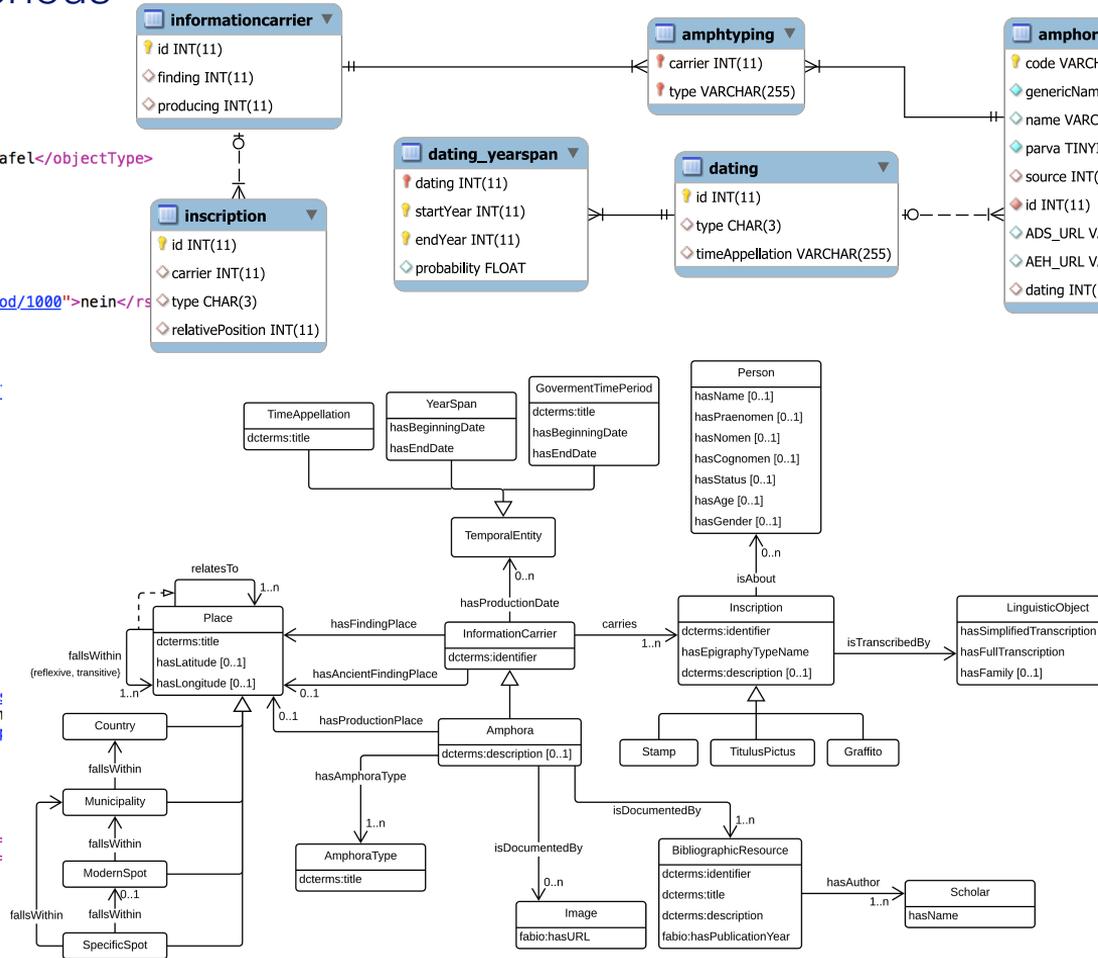
Courtesy of Museo Arqueológico de Sevilla (Dr. Fernando Fernández Gómez) Simon Keay

The available datasets are heterogenous

(both conceptually and technically)

```

<physDesc>
  <objectDesc>
    <supportDesc>
      <support>
        <objectType ref="http://www.eagle-network.eu/voc/objtyp/lo/257">Tafel</objectType>
        <material>Marmor, geädert / farbig</material>
      </support>
      <dimensions unit="cm">
        <height>33</height>
        <width>34</width>
        <depth>2.7</depth>
      </dimensions>
      <rs type="decoration" ref="http://www.eagle-network.eu/voc/decor/lo/1000">nein</rs>
    </supportDesc>
  </objectDesc>
  <handDesc>
    <handNote>
      <dimensions>
        <width unit="cm"/>
        <height unit="cm">3.2-2</height>
      </dimensions>
    </handNote>
  </handDesc>
</physDesc>
<history>
  <origin>
    <origPlace><placeName type="provinceItalicRegion" ref="http://www.trismegi:
placeName><placeName ref="http://www.trismegistos.org/place/014437">Cun
    <origDate notBefore-custom="0071" notAfter-custom="0130" datingMethod="htj
AD
    </origDate>
  </origin>
  <provenance type="found">
    <placeName ref="http://www.geonames.org/">Cuma, bei</placeName><placeName type:
www.geonames.org/">Campania</placeName><placeName type="modernCountry" ref:
placeName></provenance>
</history>
  
```





“I am looking for all the **amphoras** of type **Dressel 20** (information carriers) produced in the settlement of **Malpica**, together with all the available information about their **physical characteristics**, **dating** and carried **inscriptions**, in all the available datasets. ”

Relational model: Example

InformationCarrier

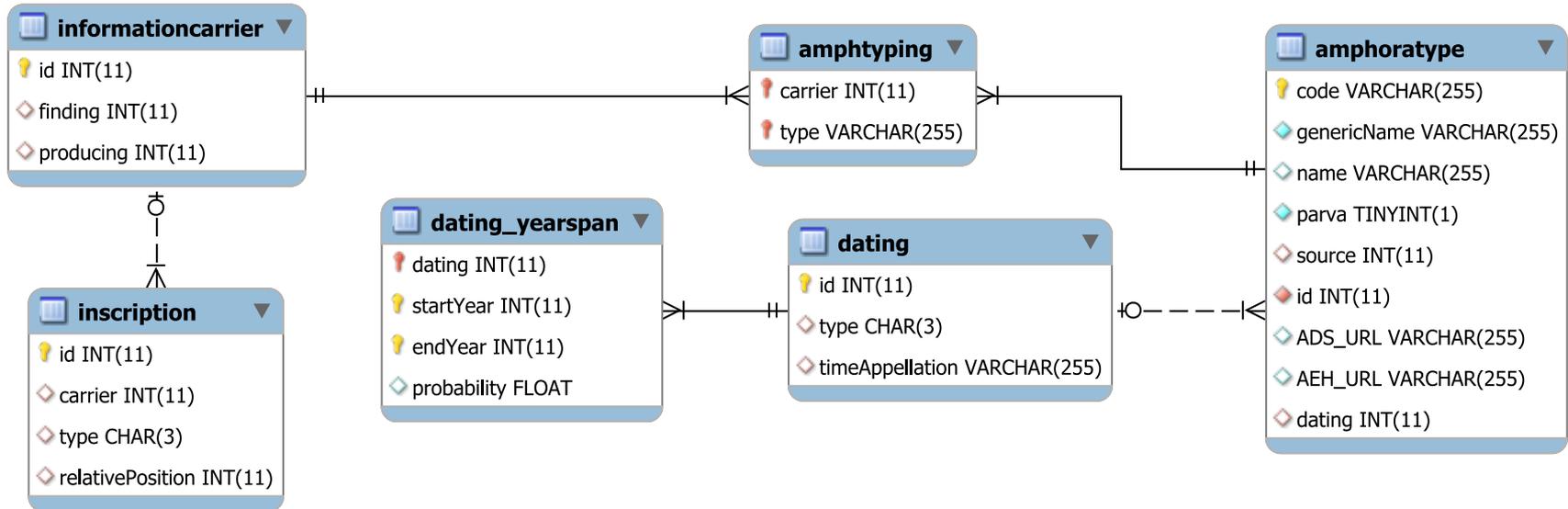
carrierId	finding	producing
C234	'Monte Testaccio'	null
C678	'Hispalis'	'Malpica'
C449	null	'Malpica'

AmphTyping

carrierId	finding
C234	'Monte Testaccio'
C678	'Hispalis'
C449	null

AmphoraType

typeCode	genericName	name	source	...
T15	Dressel	Dressel 20	null	...
T09	Africano grande	Africana 2C	S25	...



SQL queries: Example

- SQL is the standard language for querying relational databases.
- The core of SQL corresponds to relational algebra, a well-studied formalism.

	ic_id integer	d_id integer	sy integer	ey integer
1	1	1	-200	-150
2	13	7	41	54
3	17	8	216	223
4	17	8	214	214
5	84	9	216	223
6	84	9	214	214
7	198	10	216	223
8	198	10	214	214
9	212	11	216	223

“Return the amphoras and the date when they were produced”

SELECT

```
ic.id AS ic_id, d.id AS d_id, dys.startYear AS sy,  
dys.endYear AS ey
```

FROM

```
InformationCarrier ic  
JOIN Producing p ON ic.producing = p.id  
JOIN Dating d ON p.dating = d.id  
JOIN DatingYearSpan dys ON dys.dating = d.id
```

When can this go wrong?

To query this information and obtain the desired answers requires:

- To have a have deep understanding of the datasets.
- To **create the proper queries** that extract all information about each type of object from each dataset is a complex task even for experts.
- To **merge the answers** returned from each dataset, possibly filtering out undesired objects or undesired parts of the answers (e.g., the objects that were not produced in ‘Malpica’).
- To **query multiple data sources**, one can use a *data federation tool*, which exposes multiple data sources as if they were a single relational database (e.g., Teiid, Exareme) but...

Scholars can easily get lost!

The way we tackle the problem

- Exploit the knowledge that scholars have about the domain, and make this knowledge and its associated vocabulary explicit
 - ↳ **Ontology**
- Provide a good understanding of the source data by connecting it to the domain knowledge
 - ↳ **Mappings**
- Enable scholars to rapidly formulate **intuitive queries** using the ontology (which provides a familiar vocabulary and conceptualisations), and not the data sources.



Ontology-based data integration (OBDI)

Ontology

provides a unified common vocabulary, and a conceptual view of the data.

- OWL 2 QL

Mappings

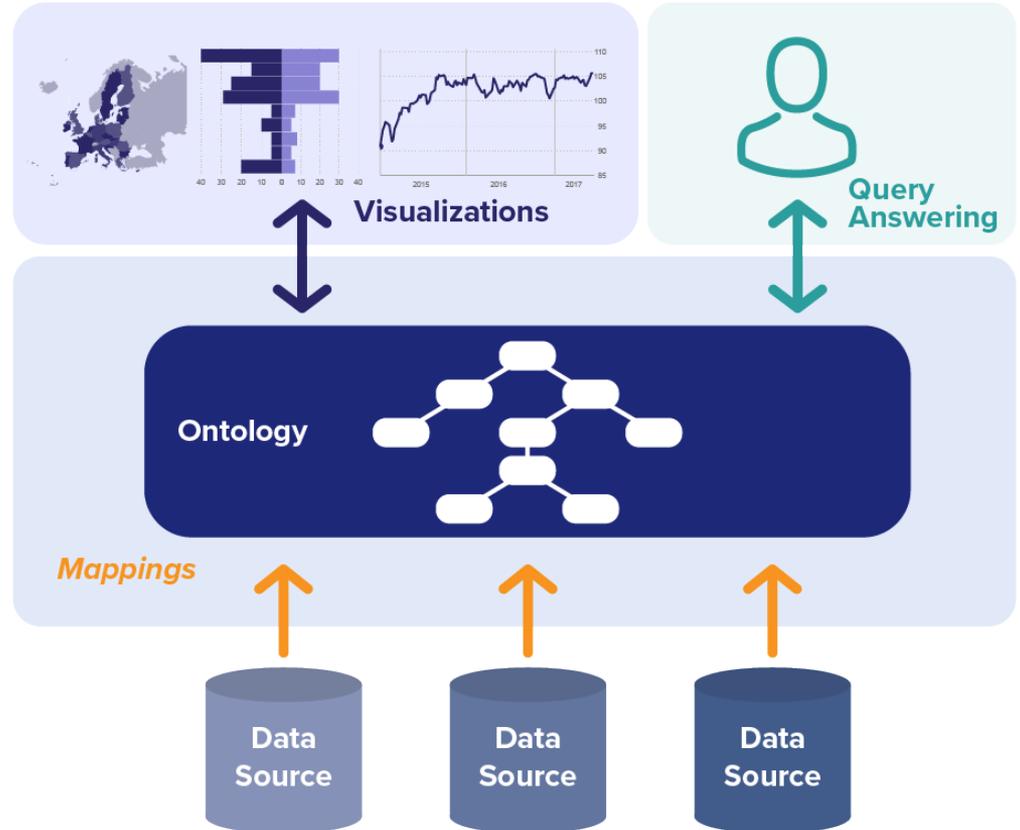
relate the terms in the ontology to the data in the sources by means of queries.

- R2RML

Data sources

are external and independent (possibly heterogeneous).

- Oracle, DB2, Postgres, MySQL, etc.



Ontologies in Computer Science

Ontologies (in Computer Science):

- are used to represent a domain of interest in a way that is comprehensible to end users
- allow for efficient processing by machines, to infer new information from the one explicitly represented.

Ontology languages:

- Are grounded in [mathematical logic](#), which makes them rigorous and not ambiguous.
- Are equipped with a [formal syntax](#), which tells us how to write expressions in the languages:
 - logic based (e.g., $\exists \text{carries} \sqsubseteq \text{Inscription}$)
 - serialised and text-based (e.g., `:carries rdf:range :Inscription`)
 - diagrammatic/graphic.
- Have a [formal semantics](#), usually provided in terms of logic.

In ontologies, the knowledge is structured into:

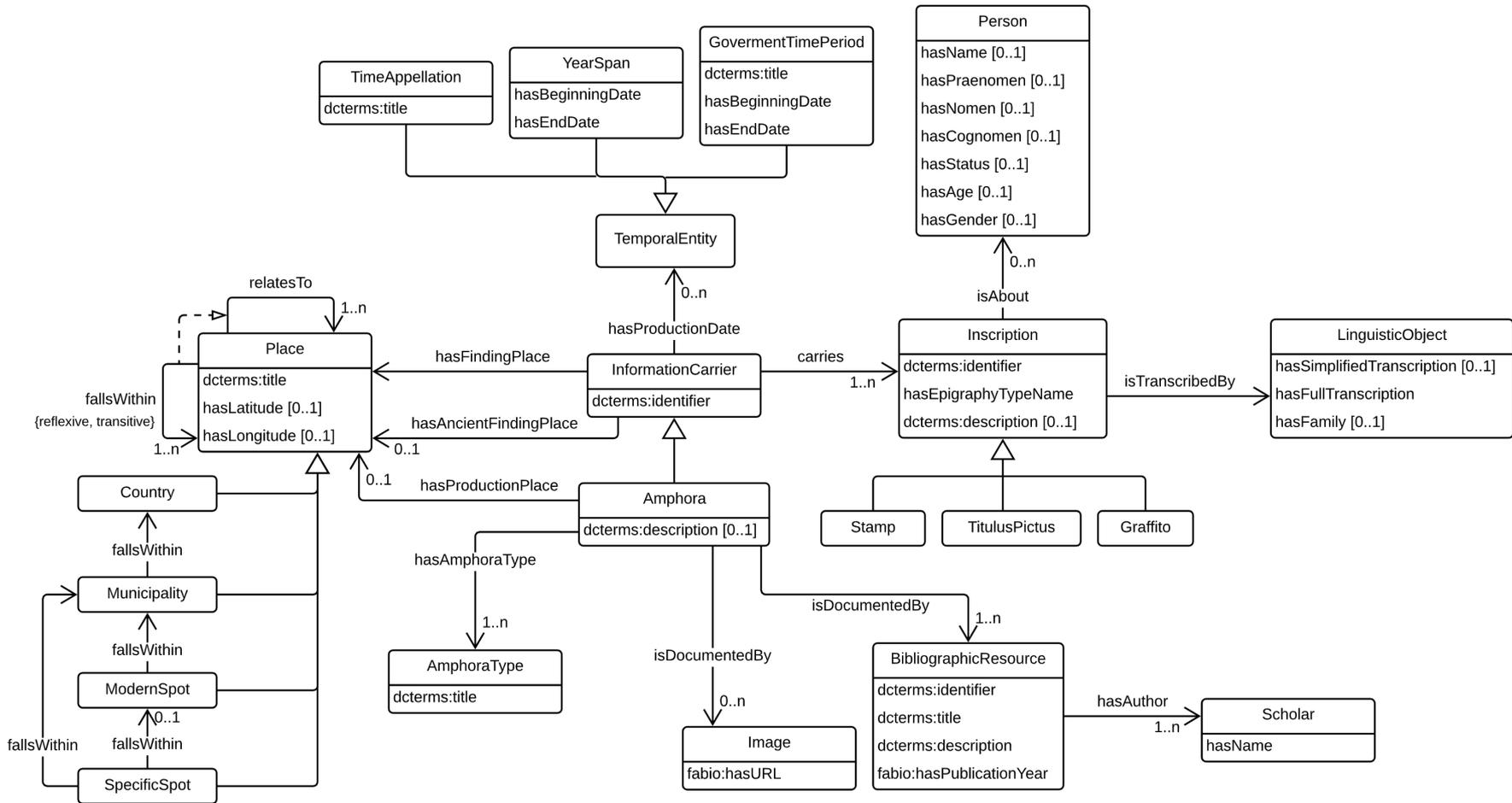
- Classes of objects (e.g., [Inscription](#), [Finding](#))
- Properties of class instances (e.g., [name](#), [startYear](#))
- Relationships between classes (e.g., [hasShapeType](#) between [Stamp](#) and [ShapeType](#))
- Properties of relation instances

The knowledge about the domain is the stated by means of (logical) [assertions](#).

Ontology development in EPNNet

In EPNNet, the OBDI team working on the ontological modelling, has made use of this correspondence:

- They have developed the **EPNet Conceptual Reference Model (CRM)**, an extensive conceptual model that *integrates* and *specialise* already existing ontologies and standards for the representation of historical data (e.g., CIDOC CRM, the EAGLE metadata model, and FaBiO).
- Using a domain-oriented vocabulary (made of terms like “inscription”, “stamp”, “simple” and “full transcription”, “grafito”, etc.), they have build a ontology expressed in the **lightweight ontology language OWL 2 QL**



Query answering over ontologies

- When a query is posed over the ontology, the OBDI system can reason over the knowledge in the ontology to provide more answers.
- This is achieved by **rewriting** the query posed by the user into a new query, that is then further processed.
- This new query incorporates the knowledge provided by the ontology (the ontology is **compiled into** the query).

Query answering by rewriting

Query answering by rewriting: Example

Query: $q(x) \leftarrow \text{worksFor}(x, y), \text{Project}(y)$

Data: $\text{worksFor}(\text{jose}, \text{epnet})$ $\text{Coordinator}(\text{jose})$
 $\text{worksFor}(\text{martin}, \text{ontop})$ $\text{Coordinator}(\text{diego})$

Query answering by rewriting: Example

Ontology:

Coordinator \sqsubseteq Researcher

Researcher $\sqsubseteq \exists$ worksFor

\exists worksFor⁻ \sqsubseteq Project

Corresponding rules:

Coordinator(x) \rightarrow Researcher(x)

Researcher(x) $\rightarrow \exists y$ (worksFor(x, y))

worksFor(x, y) \rightarrow Project(y)

Query: $q(x) \leftarrow$ worksFor(x, y), Project(y)

Perfect rewriting: $q(x) \leftarrow$ worksFor(x, y), Project(y)

$q(x) \leftarrow$ worksFor(x, y), worksFor($-, y$)

$q(x) \leftarrow$ worksFor($x, -$)

$q(x) \leftarrow$ Researcher(x)

$q(x) \leftarrow$ Coordinator(x)

Data: worksFor(jose, epnet)

worksFor(martin, ontop)

Coordinator(jose)

Coordinator(diego)

Query answering by rewriting: Example

Ontology:

Coordinator \sqsubseteq Researcher

Researcher \sqsubseteq \exists worksFor

\exists worksFor⁻ \sqsubseteq Project

Corresponding rules:

Coordinator(x) \rightarrow Researcher(x)

Researcher(x) \rightarrow $\exists y$ (worksFor(x, y))

worksFor(x, y) \rightarrow Project(y)

Query: $q(x) \leftarrow$ worksFor(x, y), Project(y)

Perfect rewriting: $q(x) \leftarrow$ worksFor(x, y), Project(y)

$q(x) \leftarrow$ worksFor(x, y), worksFor($-, y$)

$q(x) \leftarrow$ worksFor($x, -$)

$q(x) \leftarrow$ Researcher(x)

$q(x) \leftarrow$ Coordinator(x)

Data: worksFor(jose, epnet)

Coordinator(jose)

worksFor(martin, ontop)

Coordinator(diego)

Evaluating the perfect rewriting over the ABox (seen as a DB) produces as answer {jose, martin, diego}.

Query answering by rewriting: Example

Suppose we want to query the Incriptions in our database:

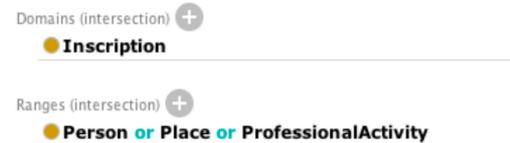
```
Select ?p WHERE { ?p rdf:type :Inscription . }
```

And our ontology says....

Subclasses of Inscription



Property :isAbout



Suppose that we only have mappings for Stamp and the property IsAbout

Without Reasoning there is no Answer!!!!

Query answering by rewriting: Example

- By looking at the mappings, we know that the original query returns no answer.
- By looking at the ontology we know that:
 - every stamp is an inscription
 - every element in the domain of `IsAbout` is an inscription.

So we can **rewrite** the original query as:

```
Select ?p WHERE {  
  { ?p rdf:type :Stamp . }  
  UNION  
  { ?p :isAbout ?y . } }
```

Mapping the data sources to the ontology

In an OBDI system, the mapping \mathcal{M} encodes how the data \mathcal{D} in the sources should be used to populate the elements of the ontology.

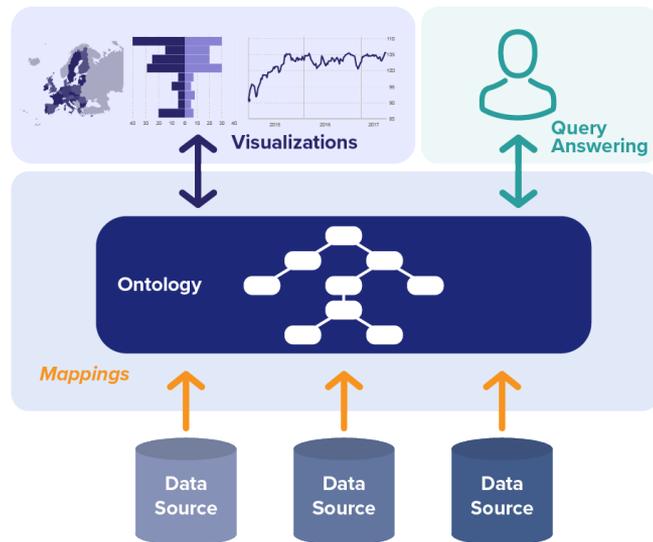
Concrete mapping languages

- Several proposals for languages to map a relational DB to an ontology have been made:
- They assume that the ontology is populated in terms of triples of the RDF data model.
- Some template mechanism is used to specify the triples to instantiate.

Virtual data layer

The data \mathcal{D} and the mapping \mathcal{M} define a virtual data layer $\mathcal{V} = \mathcal{M}(\mathcal{D})$

- Queries are answered w.r.t. the ontology and \mathcal{V} .
- We do not really materialize the data of \mathcal{V} (it is virtual!).
- Instead, the intensional information in the ontology and the mapping is used to translate queries over the ontology into SQL queries formulated over the sources.



Virtual RDF graph

The extensional counterpart of an ontology

In RDF, all data are represented by means of triples of the form: **<subject, property, object>**

- These triples form a so-called RDF graph, which is what the user actually queries.
- The mappings specify how to construct this virtual RDF graph from the data sources and the mappings .

$$t: L(t(\mathbf{x})) \leftarrow s: Q(\mathbf{x})$$

$C(t_1(\mathbf{x}_1)), p(t_1(\mathbf{x}_1), t_2(\mathbf{x}_2)),$ or $d(t_1(\mathbf{x}_1), t_2(\mathbf{x}_2)) \leftarrow$ SQL query over the DB schema

Through the mapping, each result row returned by the SQL query in the right-hand side generates a triple in the virtual RDF graph according to the triple template.

Virtual RDF graph: Example

AmphoraT

amId	place	inscription
A24	'Monte Testaccio'	'PNN'
A52	'Hispalis'	'ARVA'

Mappings:

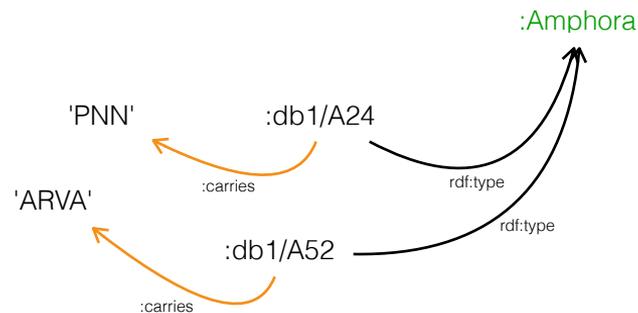
```
(:db1/{id}, rdf:type, :Amphora) ← SELECT amId AS id FROM AmphoraT
```

```
(:db1/{id}, :carries, {inscription}) ← SELECT amId AS id, inscription FROM AmphoraT
```

Virtual RDF graph: Example

AmphoraT

amId	place	inscription
A24	'Monte Testaccio'	'PNN'
A52	'Hispalis'	'ARVA'



Mappings:

`(:db1/{id}, rdf:type, :Amphora) ← SELECT amId AS id FROM AmphoraT`

generates the triples: `(:db1/A24, rdf:type, :Amphora),`
`(:db1/A52, rdf:type, :Amphora)`

`(:db1/{id}, :carries, {inscription}) ← SELECT amId AS id, inscription FROM AmphoraT`

generates the triples: `(:db1/A24, :carries, 'PNN')`
`(:db1/A52, :carries, 'ARVA')`

Using the mappings in OBDI: “Entity linking” Example

Suppose that we want now to extract the place of amphoras from a different data source:

- We need to **combine the answers** coming from different data sources.
- The different data sources might adopt **different identifiers** for data representing the same objects.

AmphoraT

amId	place	inscription
A24	'Monte Testaccio'	'PNN'
A52	'Hispalis'	'ARVA'

PlaceT

amphId	place
pl-A24	'Monte Testaccio'
pl-A52	'Hispalis'

Generating the same URI (i.e., RDF identifier):

```
(:db1/{id},:hasPlace, :db3/{place}) ← SELECT stringOp(amphId) AS id, place FROM PlaceT
```

where *stringOp* is a suitable string operation that, e.g., deletes the 'pl-' prefix from the amphora identifier.

The mappings can be used **to link the entities extracted from the different data sources**, so that at the level of the ontology they can be recognised as representing the same object.

Using the mappings in OBDI: Example

Customising the data access

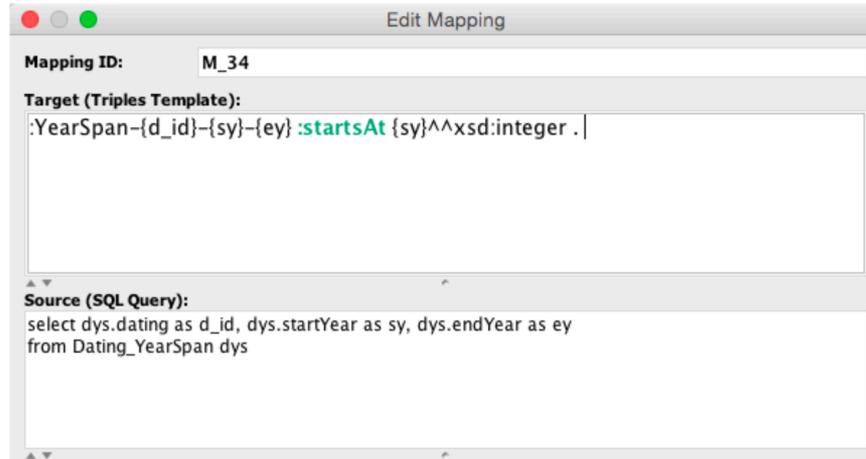
M_63 :Amphora-{ic_id}:producedAt:Tiberius-Government. select ic.id as ic_id from InformationCarrier ic join Producing n
M_64 :Amphora-{ic_id}:producedAt:Caligula-Government. select ic.id as ic_id from InformationCarrier ic join Producing n
M_65 :Amphora-{ic_id}:producedAt:Claudius-Government. select ic.id as ic_id from InformationCarrier ic join Producing n
M_66 :Amphora-{ic_id}:producedAt:Nero-Government. select ic.id as ic_id from InformationCarrier ic join Producing n

We classify the time periods by *imperial dynasties*.

- For instance, **Caligula-Government** is defined as (startYear >= 37 and endYear <= 41).

Now, we can query a government instead of integers.

Homogenising data



We homogenise all the *dates* in the different databases: integer, strings, or date, etc.

> Start a search...

Browsing 24 images so far.
Scroll down to fetch more images, if any

Current searches

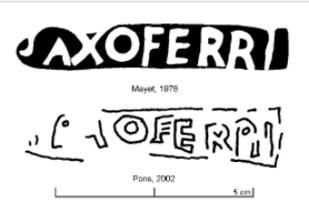
- 681 inscriptions
198 amphoras
- 2273 inscriptions
1972 amphoras
- 2022 inscriptions
558 amphoras



SimplifiedTranscription
SAXOFERREO
FullTranscription
[SAX]O'FERREO]



SimplifiedTranscription
SAXOFERREO
FullTranscription
SAXOFE'[RREO].



SimplifiedTranscription
SAXOFERRI
FullTranscription
[SAXOFERRI]



SimplifiedTranscription
SAXOFERREO
FullTranscription
[SAXI'XOFERREO]



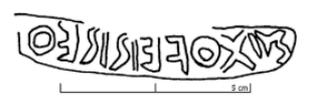
SimplifiedTranscription
SAXOFERR
FullTranscription
[S*AXOFERRI]



Ponsich (1982)



SimplifiedTranscription



SimplifiedTranscription
SAXOFERREO
FullTranscription
[S*AXOFERREO]



Timeline Faceting

AmphoricTypeTitle	AreaTitle	FindingPlaceTitle	SimplifiedTranscription
Africana 1A pic.	Aachen	(Fund-Nr.70/3)	((PALMA))
Amphora incerta	Abernethy	(PS Moseler)	...A...
August 17	Adenau	(bei Kanalbau)	...A.....
August 33	Agde	-6m; 15.2.1905	...A?...
Beltrán 2A	Alcolea del Río	-3,5m; 2.12.1939.	...AGE?...
Brindisian amphora	Alexandrie	-5-6 m.; 21.2.1905.	...AL?...
Dressel 1	Alise-Sainte-Reine	-6m; a. 14.2.1905.	...ARPL...?...
Dressel 10	Altenstadt	0004/141 (13-7-1994)	...CCL...
Dressel 2-4	Altkalkar	0004/321 (22-7-1994)	...CPRIS?...
Dressel 2-4 Baetican	Alzey	0006/503 (22-5-1995)	...DIM?...
Dressel 2-4 Catalan	Amiens	004/160	...DV....
Dressel 2-4 Loire Basin	Angers	10.11.1904	...FEK...
Dressel 20	Ardèche (Alba)	12.3.1961.	...FSS...



```

1 ► PREFIX :↔
17 SELECT ?ceipacNumber ?transcription ?findingPlace ?findingSpot ?country
18 ▼ WHERE {
19     ?amph a :Amphora .
20     ?amph dcterms:identifier ?ceipacNumber .
21     ?amph :carries ?inscription .
22     ?inscription :isTranscribedBy ?linguisticObject .
23     ?linguisticObject :hasFullTranscription ?transcription .
24     ?amph :hasFindingPlace ?findplace .
25     ?findplace :fallsWithin ?cou .
26     ?cou a :Country .
27     ?cou dcterms:title ?country .
28     FILTER (?country = "Spain")
29     ?findplace :fallsWithin ?mun .
30     ?mun a :Municipality .
31     ?mun dcterms:title ?findingPlace .
32 ▼ optional {
33     ?findplace :fallsWithin ?msp .
34     ?msp a :ModernSpot .
35     ?msp dcterms:title ?findingSpot .
36 }
37 }
38 ORDER BY ?findingPlace ?findingSpot
39 LIMIT 100
    
```

↔
Table
Response
Pivot Table
Google Chart
Geo
Execute & Download CSV

Showing 1 to 50 of 100 entries (in 2.473 seconds)

Search: Show entries

	ceipacNumber	transcription	findingPlace	findingSpot	country
1	"12585"^^xsd:int	T^H	Alayor	Alcaidús d'en Flaquer	Spain
2	"12585"^^xsd:int	F'	Alayor	Alcaidús d'en Flaquer	Spain
3	"41452"^^xsd:int	P	Alayor	Binialmesc	Spain
4	"12574"^^xsd:int	M·C·N	Alayor	Cotaina	Spain
5	"12581"^^xsd:int	PR	Alayor	Cotaina	Spain

What we have seen today

- **Ontologies** and **mappings** provide a vocabulary to formulate the queries, enrich them, and find the answers in (possibly) multiple heterogenous data sources.
- **Ontologies** and **Semantic Web Technology** can help to handle the problem of accessing and integrating data sources, also in the Cultural Heritage domain.

Diversity:

- Using ontologies describing particular domains allows to hide the *storage complexity*.
- Agreement on data identifiers allows for *integration* of multiple datasets.

Understanding:

- Agreement on a domain-oriented vocabulary allow to better define your data and allows for easy information exchange.

What we left outside

- Semantic Query Optimisation, SPARQL-based querying, Performance Evaluation, Aggregates and bag semantics, no-SQL repositories, ... *Tons of theory*.

Thank you for your attention!

 <https://romanopendata.eu/>

Questions?

